



The
University
Of
Sheffield.

MAS345

SCHOOL OF MATHEMATICS AND STATISTICS

Spring Semester 2012–2013

Codes and Cryptography

2 hours 30 minutes

Attempt all the questions. The allocation of marks is shown in brackets.

- 1 (i) Let p be prime and let C be an $[n, k]$ -linear code over \mathbb{F}_p . Show that C has precisely p^k codewords. *(3 marks)*
- (ii) Let C be the linear code over \mathbb{F}_2 that has generator matrix

$$G = \begin{pmatrix} 1 & 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{pmatrix}.$$

- (a) How many codewords are there in C ? How many codewords are there in the dual code C^\perp ? *(2 marks)*
- Let \mathcal{E}_7 be the even weight code of length 7. How many of the codewords of C are in \mathcal{E}_7 ? *(2 marks)*
- (b) Find a parity check matrix H for C . *(5 marks)*
- (iii) Let D be an $[n, k]$ -linear code over \mathbb{F}_2 and let $d \in D \cap D^\perp$. Show that $\text{wt}(d)$ is even. *(3 marks)*
- With C as in (ii), find all elements of $C \cap C^\perp$. *(3 marks)*
- (iv) Let $v = 0001110 \in \mathbb{F}_2^7$ and let S be the Hamming sphere with centre v and radius 4. How many words are there in S ? How many of these words are also in the even weight code \mathcal{E}_7 ? *(5 marks)*
- Is S a linear code? Justify your answer. *(2 marks)*

2 (i) Let p be prime and let C be an $[n, k, d]$ -linear code over \mathbb{F}_p . State, without proof, a result which expresses d in terms of linear dependence of columns of a parity check matrix for C . Deduce that $d \leq n - k + 1$ (the Singleton Bound). **(4 marks)**

(ii) Let C be an $[n, k, 3]$ -linear code over \mathbb{F}_p and let H be a parity check matrix for C with columns h_1, h_2, \dots, h_n . Let $a = a_1a_2 \dots a_i \dots a_n \in C$ and let $b = a_1a_2 \dots b_i \dots a_n$ be obtained from a by a single error which occurs in the i th bit. Show that $Hb^T = (b_i - a_i)h_i$. Show also that the position i of the error and the correct i th bit a_i are uniquely determined by Hb^T . **(5 marks)**

(iii) Let

$$H = \begin{pmatrix} 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 2 \\ 1 & 0 & 1 & 2 & 0 & 1 & 2 & 0 & 1 & 2 & 0 \end{pmatrix} \in M_{3 \times 11}(\mathbb{F}_3)$$

and let C be the linear code over \mathbb{F}_3 with parity check matrix H .

(a) Show that C has minimum distance 3. **(2 marks)**

(b) A codeword $a \in C$ is transmitted and 10101111011 is received. Explain why at least one error must have occurred and, on the assumption that a single error has occurred, identify the codeword a . **(5 marks)**

Show that 1111111111 cannot be obtained from a codeword in C by an error in a single bit but that it can be obtained from a codeword in C by errors in two bits. **(3 marks)**

(c) The columns of H are the base 3 representations, written as columns, of the first 15 positive integers with 2, 6, 7 and 8 omitted. Show that if H is replaced by a 3×12 matrix over \mathbb{F}_3 obtained from H by inserting an extra column corresponding to the base 3 representation of 2, 6, 7 or 8, then the resulting code does not correct single errors. **(2 marks)**

(d) Is C a perfect code? Justify your answer. **(4 marks)**

- 3 (i) A message of fifty-seven characters, each of which is a capital letter, is encoded in \mathbb{Z}_{26} , using the correspondence in the table on the provided data sheet, and then encrypted using the one-time pad method, deleting spaces and any punctuation marks, working modulo 26 and using the following translation of Euclid's Proposition 20 as key:
 PRIME NUMBERS ARE MORE THAN ANY ASSIGNED MULTITUDE OF PRIME NUMBERS.

The first eight characters of the encrypted message are IYQEIKUY. Decrypt these eight characters. *(3 marks)*

Later in the unencrypted message, the word EASY appears and the E is encrypted as K. How are the other three letters of EASY encrypted? *(2 marks)*

- (ii) Messages are broken up into units that can be encoded using the 27 elements of \mathbb{F}_3^3 . Elements of \mathbb{F}_3^3 are then written as columns and encrypted using 3-dimensional affine encryption modulo 3, with the transformation $f : \mathbb{F}_3^3 \rightarrow \mathbb{F}_3^3$ given by $f(V) = KV + L$, where

$$K = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 2 & 0 \\ 1 & 0 & 2 \end{pmatrix} \text{ and } L = \begin{pmatrix} 2 \\ 1 \\ 2 \end{pmatrix}.$$

Find $a, b, c, d, e \in \mathbb{F}_3$ such that K has inverse

$$\begin{pmatrix} a & 0 & b \\ 0 & c & 0 \\ d & 0 & e \end{pmatrix}.$$

(1 mark)

Find $V \in \mathbb{F}_3^3$ such that V is encrypted as $(1 \ 2 \ 0)^T$ and find $W \in \mathbb{F}_3^3$ such that W is encrypted as itself. *(6 marks)*

- (iii) (a) Gavin and Stacey are members of a group who use the ElGamal encryption method with $p = 101$ and the generator $g = 27$ of \mathbb{F}_{101} . Gavin's secret key is 69. He receives the encrypted message $(6, 6)$ from Stacey. Decrypt the message. You may use the squares in \mathbb{F}_{101} shown on the data sheet as an aid to calculation. *(5 marks)*
- (b) Horatio is a member of a consortium that uses the ElGamal encryption method with a prime number p and a generator g of \mathbb{F}_p . Let a_H be Horatio's secret key and let $(x, y) = (g^k, P(g^{a_H})^k)$ be the encryption of a message P sent to Horatio by another consortium member. Show that

$$x = y \Leftrightarrow P = x^{p-a_H}$$

and that $P = x^2 \Leftrightarrow x^{2+a_H} = y.$

(8 marks)

- 4 (i) Let p be a prime number and let the prime factorization of $p - 1$ be $p_1^{m_1} p_2^{m_2} \dots p_k^{m_k}$, where p_1, p_2, \dots, p_k are distinct prime numbers and each $m_i > 0$. For $1 \leq i \leq k$, let $q_i = (p - 1)/p_i$. Let $a \in \mathbb{F}_p^*$. Show that a is a generator of \mathbb{F}_p if and only if $a^{q_i} \neq 1$ for $1 \leq i \leq k$. **(5 marks)**

Question 3 refers to 27 as a generator of \mathbb{F}_{101} . Show that 27 is indeed a generator of \mathbb{F}_{101} . You may use the squares in \mathbb{F}_{101} shown on the data sheet. **(4 marks)**

- (ii) The superincreasing 5-tuple $(v_1 = 1, v_2 = 6, v_3, v_4, v_5)$ is chosen in such a way that each v_i is the minimal value of v_i such that $v_i \equiv 1 \pmod{5}$. Identify v_3, v_4 and v_5 . **(2 marks)**

- (iii) An organisation, to which Yvette belongs, uses Merkle-Hellman knapsack encryption, with superincreasing 5-tuples, to encrypt messages that can be expressed using the 26 letters A–Z, encoded as on the data sheet but with A encoded as 26, and five punctuation marks that are encoded as 27–31. In choosing her key, Yvette first chooses the superincreasing 5-tuple from (ii).

Her choice for an integer $M > \sum_{i=1}^5 v_i$ is 100 and her choice for an integer a , with $0 < a < M$ and $(a, M) = 1$, is 27 which, modulo 100, has inverse 63. She receives an encrypted message which, modulo 100, reduces to 53, 56, 4, 69. Identify the four-letter word that was encrypted. **(6 marks)**

- (iv) Let p_1, p_2, \dots, p_k be distinct prime numbers and let $n = p_1 p_2 \dots p_k$. State, without proof, a necessary and sufficient condition for n to be a *Carmichael number*. **(2 marks)**

Let $m > 2$ be an odd positive integer such that $6m$ divides 90 and $6m + 1$ is prime. Show that $7 \times 13 \times (6m + 1)$ is a Carmichael number. Hence find two Carmichael numbers that are divisible by 91. **(6 marks)**

End of Question Paper

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DATA SHEET

Table for Q3(i) and Q4(iii).

A	B	C	D	E	F	G	H	I	J	K	L	M
0*	1	2	3	4	5	6	7	8	9	10	11	12
N	O	P	Q	R	S	T	U	V	W	X	Y	Z
13	14	15	16	17	18	19	20	21	22	23	24	25

* In Q4(iii), A is encoded as 26.

The list below shows the nonzero squares in \mathbb{F}_{101} .

$$\begin{array}{llllllll}
 1^2 = 1, & 2^2 = 4, & 3^2 = 9, & 4^2 = 16, & 5^2 = 25, & 6^2 = 36, & 7^2 = 49, & 8^2 = 64, \\
 9^2 = 81, & 10^2 = 100, & 11^2 = 20, & 12^2 = 43, & 13^2 = 68, & 14^2 = 95, & 15^2 = 23, & 16^2 = 54, \\
 17^2 = 87, & 18^2 = 21, & 19^2 = 58, & 20^2 = 97, & 21^2 = 37, & 22^2 = 80, & 23^2 = 24, & 24^2 = 71, \\
 25^2 = 19, & 26^2 = 70, & 27^2 = 22, & 28^2 = 77, & 29^2 = 33, & 30^2 = 92, & 31^2 = 52, & 32^2 = 14, \\
 33^2 = 79, & 34^2 = 45, & 35^2 = 13, & 36^2 = 84, & 37^2 = 56, & 38^2 = 30, & 39^2 = 6, & 40^2 = 85, \\
 41^2 = 65, & 42^2 = 47, & 43^2 = 31, & 44^2 = 17, & 45^2 = 5, & 46^2 = 96, & 47^2 = 88, & 48^2 = 82, \\
 49^2 = 78, & 50^2 = 76, & 51^2 = 76, & 52^2 = 78, & 53^2 = 82, & 54^2 = 88, & 55^2 = 96, & 56^2 = 5, \\
 57^2 = 17, & 58^2 = 31, & 59^2 = 47, & 60^2 = 65, & 61^2 = 85, & 62^2 = 6, & 63^2 = 30, & 64^2 = 56, \\
 65^2 = 84, & 66^2 = 13, & 67^2 = 45, & 68^2 = 79, & 69^2 = 14, & 70^2 = 52, & 71^2 = 92, & 72^2 = 33, \\
 73^2 = 77, & 74^2 = 22, & 75^2 = 70, & 76^2 = 19, & 77^2 = 71, & 78^2 = 24, & 79^2 = 80, & 80^2 = 37, \\
 81^2 = 97, & 82^2 = 58, & 83^2 = 21, & 84^2 = 87, & 85^2 = 54, & 86^2 = 23, & 87^2 = 95, & 88^2 = 68, \\
 89^2 = 43, & 90^2 = 20, & 91^2 = 100, & 92^2 = 81, & 93^2 = 64, & 94^2 = 49, & 95^2 = 36, & 96^2 = 25, \\
 97^2 = 16, & 98^2 = 9, & 99^2 = 4, & 100^2 = 1.
 \end{array}$$